

# Lecture 6

## Last time

- ▷ Gordon's Theorem of alternatives
- ▷ Optimality conditions with functional constraints.

## Today

- ▷ Fenchel conjugate
- ▷ Fenchel duality

## Fenchel conjugate

We will introduce a transformation that plays a key role in duality.

Question: Given a function  $f: E \rightarrow \bar{\mathbb{R}}$  and a point  $z \in E$ , how to find  $\bar{x}$  s.t.  $z \in \partial f(\bar{x})$ ?

Recall that  $g \in \partial f(\bar{x})$  iff

$$f(\bar{x}) - \langle z, \bar{x} \rangle \leq f(x) - \langle z, x \rangle \quad \forall x \in E.$$

This is equivalent to

$$\bar{x} \in \operatorname{argmin} f(x) - \langle z, x \rangle$$

$$\Downarrow$$
$$\bar{x} \in \operatorname{argmax} \langle z, x \rangle - f(x).$$

This is the core idea behind the following definition.

**Def:** The Fenchel conjugate of a function  $f: E \rightarrow \bar{\mathbb{R}}$  is

$$f^*(z) = \sup_{x \in E} \langle z, x \rangle - f(x).$$

→

**Lemma:**  $f^*$  is convex and closed.

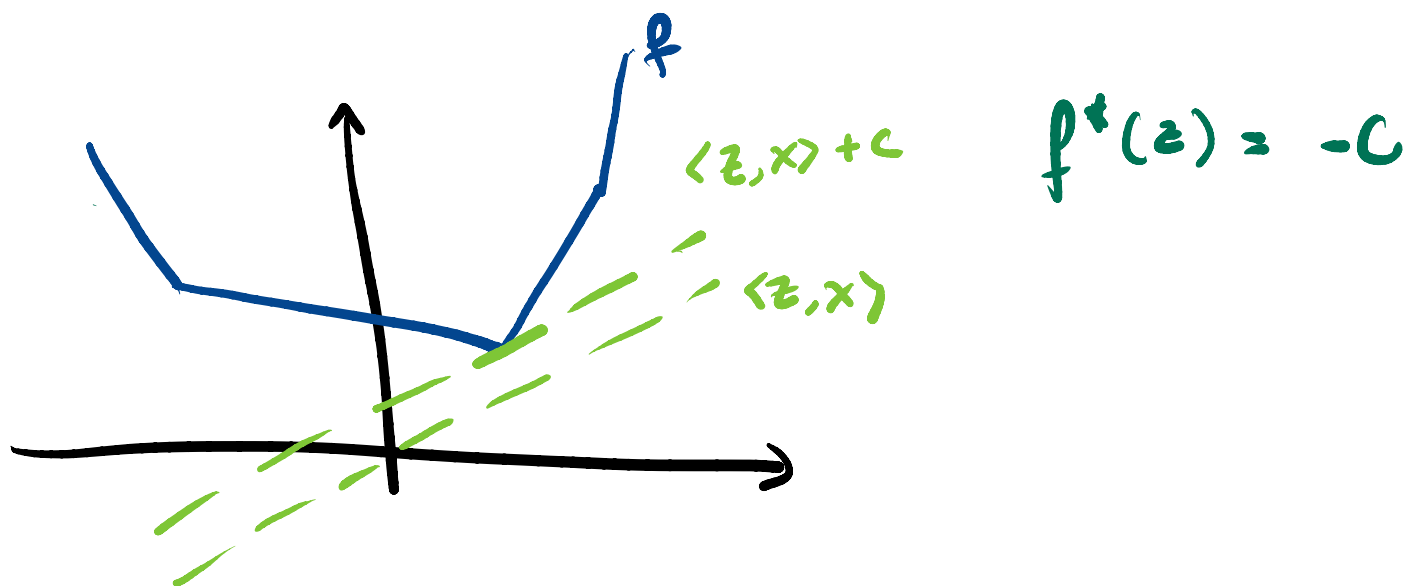
**Proof:** Note that

$$t \geq f^*(z) \iff t \geq \langle z, x \rangle - f(x) \quad \forall x \in E.$$

Then,  $\text{epi } f^*$  is an intersection of halfspaces.  $\square$

Another interpretation

By our derivation, the sup in  $f^*(z)$  is attained whenever  $z \in \partial f(x)$ . Thus, we could compute  $f^*$  geometrically by raising the graph of  $x \mapsto \langle z, x \rangle + c$  until it is tangent to  $f$ .



**Proposition:** For any  $f: E \rightarrow \bar{\mathbb{R}}$  we have that

$$f(x) + f^*(z) \geq \langle z, x \rangle$$

If  $f(x)$  is finite, then equality holds iff  $z \in \partial f(x)$ .

Proof: Follows easily by our discussion above.  $\square$

## Examples

▷ Linear functions      Suppose  $f(x) = \langle c, x \rangle$

Then,

$$\begin{aligned} f^*(z) &= \sup_x \langle z, x \rangle - \langle c, x \rangle \\ &= \sup_x \langle z, -c, x \rangle \\ &= \begin{cases} 0 & \text{if } z = c \\ +\infty & \text{otherwise.} \end{cases} \\ &= \tau_{\{c\}}(z). \end{aligned}$$

▷ Indicator of a point       $f(x) = \tau_{\{c\}}(x)$

$$\begin{aligned} f^*(z) &= \sup_x \langle z, x \rangle - \tau_{\{c\}}(x) \\ &= \sup_{x \in \{c\}} \langle z, x \rangle \\ &= \langle c, z \rangle. \end{aligned}$$

← We got  $f = f^{**}$ . We will see this is a recurrent pattern for convex closed functions.

▷ Indicator of  $[-1, 1]$  suppose  $f(x) = \chi_{[-1, 1]}(x)$

$$\begin{aligned} f^*(z) &= \sup_x z x - \chi_{[-1, 1]}(x) \\ &= \sup_{x \in [-1, 1]} z x \\ &= \text{sign}(z) z \\ &= |z|. \end{aligned}$$

▷ Absolute value  $f(x) = |x|$

$$f^*(z) = \sup_x z x - |x|$$

Let's consider two cases if  $|z| \leq 1$

$$\Rightarrow \sup_x x (z - \text{sign}(x)) \leq 0 \quad \leftarrow \text{achieved with } x=0$$

↑  
always have opposite signs

if  $|z| > 1$ , then taking any  $x = \lambda z$  for any  $\lambda > 0$  yields

$$f^*(z) \geq \lambda z^2 - \lambda |z| > \lambda (z^2 - z) > 0$$

Taking  $\lambda \uparrow \infty$  gives  $f^*(z) = \infty$ .

Therefore,  $f^*(z) = \chi_{[-1, 1]}(z)$ .

Once more the two coincide.

▷  $p$  norms

Suppose that  $f(x) = \frac{|x|^p}{p}$   $p \geq 1$ .

Then

$$f^*(z) = \frac{|z|^q}{q} \quad \text{with} \quad \frac{1}{p} + \frac{1}{q} = 1.$$

(Exercise)

# Fenchel duality

Recall that the adjoint  $A^*$  of a linear map  $A: E \rightarrow Y$  is defined via

$$\langle Ax, y \rangle = \langle x, A^*y \rangle \quad \forall x \in E, y \in Y.$$

**Theorem (Fenchel Duality):** For any functions  $f: E \rightarrow \bar{\mathbb{R}}$  and  $g: Y \rightarrow \bar{\mathbb{R}}$ , and linear map  $A: E \rightarrow Y$ . Define the primal problem

$$p^* = \inf_{x \in E} f(x) + g(Ax)$$

and the dual problem

$$d^* = \sup_{z \in Y} -f^*(A^*z) - g^*(-z).$$

Then, without extra assumptions,

$$p^* \geq d^* \quad (\text{Weak duality}).$$

If further  $f$  and  $g$  are never  $-\infty$  and convex, and

$0 \in \text{int}(\text{dom } g - A \text{dom } f) := \text{int}\{u - Av \mid u \in \text{dom } g, v \in \text{dom } f\}$   
(constraint qualification)

then, we have

$$p^* = d^* \quad (\text{Strong duality}).$$

In this case  $d^*$  is finite, and it is always attained. Moreover,  $\bar{x}$  and  $\bar{y}$  are optimal if, and only if,

$$A^* \bar{y} \in \partial f(\bar{x}) \quad \text{and} \quad -\bar{y} \in \partial g(A\bar{x})$$

(Complementary slackness)

Moreover, constraint qualification holds if either

1. There is  $\hat{x} \in \text{dom } f$  s.t.  $A\hat{x} \in \text{int dom } g$ .
2. There is  $\hat{x} \in \text{int dom } f$  s.t.  $A\hat{x} \in \text{dom } g$ , and  $A$  is surjective.

Proof: Weak duality follows from Fenchel-Young since  $\forall (x, y) \in E \times Y$

$$\begin{aligned} f(x) + f^*(A^*y) &\geq \langle x, A^*y \rangle \\ + \quad g(Ax) + g^*(-y) &\geq \langle Ax, -y \rangle \end{aligned} \quad (\text{smiley})$$

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$$f(x) + g(Ax) + f^*(A^*y) + g^*(y) \geq 0$$

Taking an inf yields  $p^* \geq d^*$ .

The key object to prove strong duality is the value function

$$v(z) := \inf_{x \in E} \{ f(x) + g(Ax + z) \} \quad \text{for } z \in Y.$$

So  $p^* = v(0)$ . Note that if  $p = -\infty$ , we are strong duality holds trivially by weak duality. Assume that  $z \in \text{dom } v$ , thus the infimum is not  $+\infty$ . Therefore,  $z \in \text{dom } v$  iff  $x \in \text{dom } f(x)$  and  $Ax + z \in \text{dom } g$ .

Thus,  $z \in \text{dom } g - A \text{dom } f$ . Hence, constraint qualification is equivalent to  $0 \in \text{int dom } v$ .

Claim (Exercise): The function  $v$  is convex.

Then, since  $0 \in \text{int dom } v$ , we have  $v(0)$  is finite and by HW1 P1c,  $v$  is never  $-\infty$  in  $\text{int dom } v$ . Thus, there exists  $-y \in \partial v(0)$ .

Moreover,

$$v(0) + v^*(-y) = \langle 0, -y \rangle = 0.$$

Therefore,

$$\begin{aligned} d^* &\leq p^* = v(0) \\ &= -v^*(-y) \\ &= -\sup_z \langle y, z \rangle - v(z) \\ &= \inf_z v(z) + \langle y, z \rangle \\ &= \inf_z \inf_x \{ f(x) + g(Ax + z) \} + \langle y, z \rangle \\ &= \inf_x \underbrace{\inf_z \{ \langle y, Ax + z \rangle + g(Ax + z) \}}_{-g^*(-y)} + \langle y, -Ax \rangle + f(x) \end{aligned}$$

$$\begin{aligned}
 &= -g^*(-y) + \overbrace{\inf \langle A^*y, x \rangle}^{-f^*(A^*y)} + f(x) \\
 &= -g^*(-y) - f^*(A^*y) \\
 &\leq d^*
 \end{aligned}$$

Thus, by weak duality  $p^* = -g^*(y) - f^*(A^*y) = d^*$ . Complementary slackness ensures that the inequalities in (:) become equalities.

Now, let's prove that constraint qualification holds given (1) or (2). If (1) holds

Then,  $0 \in \text{dom } g - A \text{dom } f$  and since  $A\hat{x} \in \text{int dom } g$   $\exists r > 0$  s.t.  $0 \in A\hat{x} + rB - A\hat{x}$

$\subseteq \text{dom } g - A \text{dom } f$ , thus  $0 \in \text{int}(\text{dom } g - A \text{dom } f)$ .

If (2) holds, then we can use the following.

**Theorem (Open Mapping)** Any surjective mapping  $A: E \rightarrow Y$  maps open sets to open sets.

We have that  $\exists r > 0$  s.t.  $\hat{x} + rB \subseteq \text{dom } f$ .

Therefore  $A\hat{x} + r\hat{A}B$  is an open map contained in  $A \text{dom } f$ . A similar argument as before, yields that constraint qualification holds.  $\square$



## Proof of the Open Mapping Theorem: A is one-to-one

We start by assuming  $\dim Y = \dim E$ .

Notice it suffices to show  $0 \in \text{int } AB$ .

Seeking contradiction assume  $0 \notin \text{int } AB$ ,

$\Rightarrow \exists (x_n) \in B^c$  s.t.  $(Ax_n)_n \in (AB)^c$  with

$Ax_n \rightarrow 0$ . Therefore,  $\frac{Ax_n}{\|x_n\|} \rightarrow 0$ .

WLOG assume  $\frac{x_n}{\|x_n\|} \rightarrow x^*$ . This is a contradiction since  $Ax^* \neq 0$ .

If  $\dim E > \dim Y$ , we reduce to the previous case. Again assume  $0 \notin \text{int } AB$ . From Linear Algebra we know

$$A(\text{Ker } A)^\perp = \text{range } A \text{ and } A(B \cap (\text{Ker } A)^\perp) = AB.$$

We can run the same argument setting

$$E' = (\text{Ker } A)^\perp.$$

□