

# Lecture 1

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Office Hours: Th 3:00 - 4:30 pm  
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## TAs:

- Thabo Samakhona (tsamakh1@jhu.edu)

OH: - am

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OH: Tu 9:15 - 10:00 pm

## Resources

- Canvas

- website (mateodd25.github.io/

- Piazza

nonlinear2

- Gradescope



Ask your questions here

Delete this 2 for resources from Nonlinear 1.

All submissions

# Agenda

- ▷ Syllabus
- ▷ Motivation
- ▷ Overview

## Syllabus

Four components:

- Homework
- Midterm Takehome ( - )
- Final Takehome ( - )
- Participation
- Engaging in class, OH, Piazza.

3 q's about theory  
1 code question  
(Python please)

Might change

## Grading System

Let  $C_H, C_M, C_F, C_P$  denote your  
normalize grades (0 - 1).

Let  $H, M, F$  be variable weights for each component.

Your grade will be the optimal value of

$$(C_H - C_P)H + (C_M - C_P)M + (C_F - C_P)F$$

$$\max C_H \cdot H + C_M \cdot M + C_F \cdot F + C_P \cdot (1 - H - M - F)$$

s.t.

$$(H, M, F) \in \mathbb{R}^3$$

$$H + M + F \leq 100$$

$$H, M \geq 15$$

$$F \geq M$$

$$M + F \leq 80$$

$$M + F \geq 50$$

$$H + M + F \geq 90$$

(P)

$p^* =$

## Textbook

We will not follow any particular textbook. We will see website for suggested references.

# Motivation

The goal of this class is to study problems of the form

$$\min_{x \in C} f(x)$$

unlike in nonlinear we will assume  $C \subseteq \mathbb{R}^d$ . This adds additional complications but gives rise to beautiful theory.

We consider two types of constrained sets:

▷ Inequality constraints

$$C = \{x \in \mathbb{R}^d \mid g_i(x) \leq 0\}$$

with  $g_i: \mathbb{R}^d \rightarrow \mathbb{R}$  differentiable.

▷ Structured convex

We let  $C$  be a "structured" convex set. structured comes in different flavors, i.e.,

a) We can project to  $C$ .

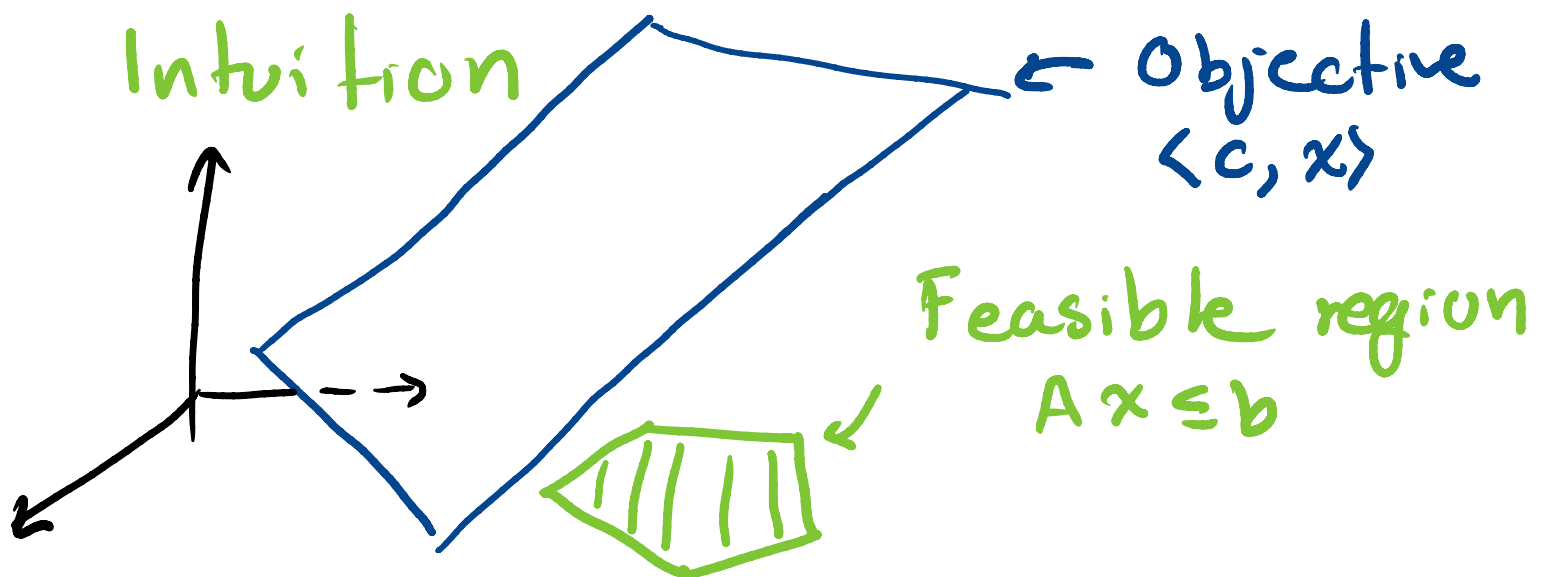
b) We have a "nice" description of  $C$ .

## Examples

▶ The grading scheme  
The problem we use to compute grades (P) is part of an important problem class called Linear Programming problems:

$$\min_{x \in \mathbb{R}^d} \langle c, x \rangle \quad \leftarrow c^T x = \sum c_i x_i$$
$$Ax \leq b \quad \leftarrow \text{component-wise.}$$

$A \in \mathbb{R}^{m \times d}$



**History Aside:** These problems and algorithms to solve them were one of the first applications of computers in the '40s.

**Question:** How can we certify that a solution is optimal?

Let's consider two students

▷ **Student 1:**

Suppose Arisu obtained

$$(C_H, C_M, C_F, C_P) = (0, 0, 100, 0) / 100$$

Bad Arisu ☹️

I claim that the grade Arisu should get is

65 out of 100  
given by weights

$$(H, M, F, P) = (20, 15, 65, 0)$$

How can I show this is the best?

$$\begin{aligned} & 1 \cdot (M + F \leq 80) \\ & + -1 \cdot (M \geq 15) \end{aligned}$$

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For every rubric  $\Rightarrow F \leq 65$

$$\langle (H, M, F, P), (C_H, C_M, C_F, C_P) \rangle$$

▷ Student 2 <sup>↑ Normalized</sup>

Suppose Boris got

$$(C_H, C_M, C_F, C_P) = (50, 50, 100, 0)_{/100}$$

Seems like nobody participates!

I claim Boris would get

82.5 out of 100 by (20, 15, 65, 0)

How to certify this?

$$\begin{aligned} & \frac{1}{2} \cdot (H + M + F \leq 100) \\ & - \frac{1}{2} \cdot (M \geq 15) \\ & + \frac{1}{2} \cdot (M + F \leq 80) \end{aligned}$$

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$$\frac{1}{2} H + \frac{1}{2} M + F \leq 82.5$$

← once again for every rubric.

It seems like this goes beyond these two students, we could have  $\lambda_1, \dots, \lambda_7 \in \mathbb{R}$

$$\lambda_1 \cdot (H + M + F \leq 100)$$

$$\lambda_2 \cdot (H \geq 15)$$

$$\lambda_3 \cdot (M \geq 15)$$

$$\lambda_4 \cdot (-M + F \geq 0)$$

$$\lambda_5 \cdot (M + F \geq 50)$$

$$\lambda_6 \cdot (M + F \leq 80)$$

$$+ \lambda_7 \cdot (H + M + F \geq 90)$$

$$(\lambda_1 + \lambda_2 + \lambda_7) \cdot H$$

$$+ (\lambda_1 - \lambda_4 + \lambda_5 + \lambda_6 + \lambda_7) \cdot M =$$

$$+ (\lambda_1 + \lambda_3 + \lambda_4 + \lambda_5 + \lambda_6 + \lambda_7) F$$

We want

$$(C_H - C_P) H$$

$$+ (C_M - C_P) M$$

$$+ (C_F - C_P) F$$

$$\leq 100 \lambda_1 + 15 \lambda_2 + 15 \lambda_3 + 50 \lambda_5 + 80 \lambda_6$$

$$+ 90 \lambda_7$$



In order to have a valid bound we need two things:

a) The coefficients in front of H, M and F have to match the grades.

b) The  $\lambda$ 's have to satisfy  
For each " $\leq$ " constraint we need  $\lambda_i \geq 0$   
For each " $\geq$ " " " " "  $\lambda_i \leq 0$ .

This leads to the problem

$$\min 100 \lambda_1 + 15 \lambda_2 + 15 \lambda_3 + 80 \lambda_5 + 80 \lambda_6 + 90 \lambda_7$$

$$\text{s.t. } \lambda_1 + \lambda_2 + \lambda_7 = C_H - C_P$$

$$\lambda_1 - \lambda_4 + \lambda_5 + \lambda_6 + \lambda_7 = C_M - C_P$$

$$\lambda_1 + \lambda_3 + \lambda_4 + \lambda_5 + \lambda_6 + \lambda_7 = C_F - C_P$$

$$\lambda_1 \geq 0, \lambda_2 \leq 0, \lambda_3 \leq 0, \lambda_4 \leq 0, \lambda_5 \geq 0$$

$$\lambda_6 \leq 0, \lambda_7 \leq 0.$$

This is another LP called

## the dual!

We have established that the  $p^* \leq d^*$ . But are they equal? Indeed, they are; this is known as "strong duality."

## Overview

We will cover:

- ▷ Fundamentals
  - ▷ Background in convexity
  - ▷ Optimality conditions
- ▷ Duality
  - ▷ Lagrange duality
  - ▷ Fenchel duality
- ▷ Algorithms
  - ▷ Classical methods
  - ▷ Splitting methods
  - ▷ Other first order methods

- ▷ Variational Analysis
  - ▷ Nonconvex calculus
  - ▷ Inverse problems & metric regularity
- Time permitting.