

Lecture 22

Last time

- ▷ Radamacher com
plexity
- ▷ Polynomial discri
mination

Today

- ▷ Vapnik - Chervonenkis
(VC) Theory.

VC Theory

We will develop another method to certify polynomial discrimination of $\{0, 1\}$ -valued functions $\mathcal{F}$.

Def: We say that

$$\mathcal{X} = \{x_1, \dots, x_n\}$$

is shattered by $\mathcal{F}$ if

$$\# \mathcal{F}(\mathcal{X}) = 2^{\#\mathcal{X}}$$

The VC dimension is

$$vc(\mathcal{F}) = \sup \{n \in \mathbb{N} \mid \exists x \in \mathcal{X}^n \text{ shattered by } \mathcal{F}\}.$$

$\uparrow$
with

We will use the following notation, if S is a class of set and

$$\mathcal{F} = \{1_A \mid A \in S\},$$

then,

$$S(X) = \mathcal{F}(X) \text{ and } VC(S) = VC(\mathcal{F}).$$

Example: Recall the last example from last class

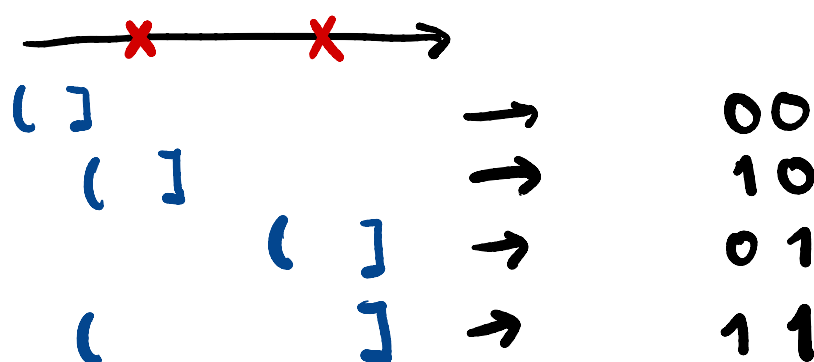
$$S_{\text{one}} = \{(-\infty, t] \mid t \in \mathbb{R}\}$$

then, we have that $VC(S_{\text{one}}) = 1$.

Similarly,

$$S_{\text{two}} = \{(a, b] \mid a, b \in \mathbb{R}, a < b\}.$$

With $n=2$



But with $n=3$ we cannot form
101.

Lemma (Sauer and Shelah): Suppose that $d = VC(S)$. Then for any $X = \{x_1, \dots, x_n\}$,

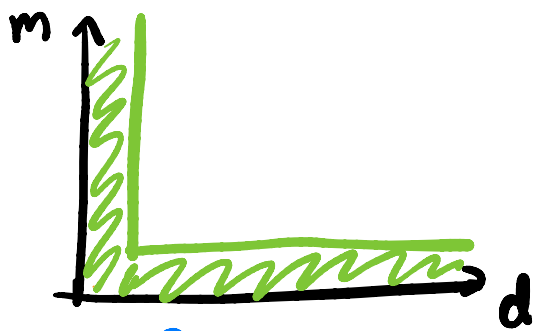
$$\# S(X) \stackrel{(1)}{\leq} \sum_{i=0}^d \binom{n}{i} \stackrel{(2)}{\leq} (n+1)^{VC(S)}.$$

$\Phi_d(n)$ $\binom{n}{k} = \begin{cases} \frac{n!}{k!(n-k)!} & 0 \leq k \leq n \\ 0 & \text{else} \end{cases}$

Thus, $\mathcal{F}$ has polynomial discrimination of order $VC(S)$ and by the results from last lecture we can bound

$$R_n(\mathcal{F}) \leq \sqrt{\frac{2VC(S) \log(n+1)}{n}}.$$

Proof: We prove (1) and leave (2) as an exercise. We will prove (1) using induction on $d+n$



Base case



Inductive step

Base case: Assume d arbitrary and $n=0$,

$$\# \mathcal{F}(\emptyset) = 0 \leq 1 = \sum_{i=0}^d \binom{0}{i} = \Phi_d(0)$$

$\binom{0}{0} = 1$

Assume n arbitrary $d=0$.

$$\# \mathcal{F}(X) = 1 \leq 1 = \binom{n}{0} = \Phi_0(n).$$

Inductive step: Fix d and m and suppose the inequality holds for any d', m' s.t. $d' + m' \leq d + m$. Let X be such that

$$X = \underset{|X|=m}{\operatorname{argmax}} \# \mathcal{F}(X).$$

Take any $a \in X$, and define

$$X^- = X \setminus \{a\}.$$

Label $H = \mathcal{F}(X) \subseteq \{0, 1\}^{\#S}$, and define

$$H_0 = \{h|_{X^-} \mid h \in H \text{ and } h(a) = 0\},$$

$$H_1 = \{h|_{X^-} \mid h \in H \text{ and } h(a) = 1\}.$$

Define $H_n = H_0 \wedge H_1$ these are binary strings that can be extended to X with both labels.

Example

$$H = \left\{ \begin{array}{ccc} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 1 \end{array} \right\} \Rightarrow H_0 = \{000\}, \quad H_1 = \{101\} \\ \text{and } H_n = \{101\}$$

Therefore,

$$\# H = \#(H_0 \cup H_1) + \# H_n.$$

Convince yourself of this!

$$\begin{aligned} &= \# \mathcal{F}(X^-) + \# H_n \quad (*) \\ \text{Inductive hypothesis} \Rightarrow &\leq \Phi_d(n-1) + \# H_n. \end{aligned}$$

Consider now

$$\mathcal{F} = \left\{ g: X^- \rightarrow \{0,1\} \mid \begin{array}{l} \exists f_1, f_2 \in \mathcal{F} \text{ s.t. } f_1(a) \neq f_2(a) \\ \text{and } \forall x \in X^-, g(x) = f_1(x) = f_2(x) \end{array} \right\}$$

Then

$$H_n \subseteq \mathcal{F}_a(X^-).$$

We claim that $vc(\mathcal{F}_a) \leq d-1$. Suppose that it is not $\Rightarrow \mathcal{F}_a$ can shatter some $Y \subseteq X^-$ and using the two extensions at a we would have

$$\# \mathcal{F}(Y \cup \{a\}) = 2^{d+1}$$

$\Downarrow$

Thus, we obtain

$$\# H_n \leq \# \mathcal{F}_a(X^-) \stackrel{\text{Inductive hypothesis}}{\leq} \Phi_{d-1}(n-1). \quad (B)$$

Fact (Pascal's Identity):

$$\binom{n}{i} = \binom{n-1}{i} + \binom{n-1}{i-1} \quad (\text{!})$$

Number of [↑] subsets of $[n]$ of size i

$$= \# \{ S \subseteq [n-1] \mid \# S = i \}$$

$$+ \# \{ S \cup \{n\} \mid S \subseteq [n-1], \# S = i-1 \}$$

Then, $(*)$, $(\heartsuit)$ and (!)

$$\begin{aligned} \# \mathcal{F}(X) &\leq \sum_{i=0}^d \binom{n-1}{i} + \sum_{i=0}^{d-1} \binom{n-1}{i} \\ &= \sum_{i=0}^d \binom{n-1}{i} + \binom{n-1}{i-1} \\ &= \sum_{i=0}^d \binom{n}{i}. \end{aligned}$$

□